INSTRUCTIONS: Solve three out of four questions. You do not have to prove results which you rely upon, just state them clearly.

## Good luck!

Q1) Solve (a), (b), (c), (d), (e).
(a) Define the $n \times n$ Vandermonde matrix $V_{n}$ (with the nodes $x_{1}, x_{2}, \ldots, x_{n}$ ), and derive the factorization:

$$
V_{n}=\underbrace{\left[\begin{array}{ccccc}
1 & 0 & 0 & \cdots & 0 \\
1 & 1 & 0 & & \vdots \\
1 & 0 & 1 & \ddots & \vdots \\
\vdots & \vdots & \ddots & \ddots & 0 \\
1 & 0 & \cdots & 0 & 1
\end{array}\right]\left[\begin{array}{ccccc}
1 & 0 & \cdots & \cdots & 0 \\
0 & x_{2}-x_{1} & \ddots & & \vdots \\
\vdots & \ddots & x_{3}-x_{1} & \ddots & \vdots \\
\vdots & & \ddots & \ddots & 0 \\
0 & \cdots & \cdots & 0 & x_{n}-x_{1}
\end{array}\right]}_{L_{1}^{-1}} \quad \underbrace{\left[\begin{array}{ccc}
1 & 0 \\
0 & V_{n-1}
\end{array}\right]}_{U_{1}^{-1}} \begin{array}{ccccc}
{\left[\begin{array}{ccccc}
1 & x_{1} & x_{1}^{2} & \cdots & x_{1}^{n-1} \\
0 & 1 & x_{1} & \ddots & \vdots \\
\vdots & \ddots & 1 & \ddots & x_{1}^{2} \\
\vdots & & \ddots & \ddots & x_{1} \\
0 & \cdots & \cdots & 0 & 1
\end{array}\right]}
\end{array}
$$

(b) Derive the formula for the determinant of $V_{n}$. Use the condition

$$
x_{i} \neq x_{j} \quad \text { for } \quad i \neq j,
$$

to prove that the Vandermonde matrix is nonsingular.
(c) Use (b) to prove that the following classical interpolation problem has a unique solution.

- Given $n$ support points

$$
\left(x_{i}, f_{i}\right) \quad i=1, \ldots, n ; \quad\left(x_{i} \neq x_{j} \quad \text { for } \quad i \neq j\right)
$$

- Find a polynomial $P(x)$ whose degree does not exceed $(n-1)$ such that

$$
P\left(x_{i}\right)=f_{i}, \quad i=1, \ldots n .
$$

(d) Use (a) to recursively to derive the formula for factoring $V_{n}^{-1}$ into a product of $n-1$ lower triangular matrices and $n-1$ upper triangular matrices. Use it to derive the Bjorck-Pereyra algorithm for solving the interpolation problem of (c).
(e) Prove that the Bjorck-Pereyra algorithm has the cost of $O\left(n^{2}\right)$ operations

Q2) Answer 3 out of 4 questions (a), (b), (c), (d).
(a) Let $\|x\|$ denotes the usual Euclidean norm $\sqrt{x^{T} x}$. Prove that the linear least squares problem

$$
\min _{x \in \mathbb{R}^{n}}\|y-A x\|
$$

with a $m \times n$ matrix $A$ has at least one minimal point $x_{0}$.
(b) Prove that if $x_{1}$ is another minimum point, then $A x_{0}=A x_{1}$. The residual $r:=y-A x$ is uniquely determined and satisfies the equation $A^{T} r=0$.
(c) Prove that Every minimum point $x_{0}$ is also a solution of normal equations

$$
A^{T} A x=A^{T} y
$$

and conversely.
(d) Explain how the orthogonalization technique (that is, computing for the $m \times n$ matrix $A$ the factorization $A=Q R$ with $m \times m$ orthogonal matrix $Q$ and $m \times n$ upper triangular matrix $R$ ) yields an efficient algorithm for solving the above least squares problem.

Q3) Answer 3 out of 4 questions (a), (b), (c), (d).
(a) Let $T$ be an $n \times n$ positive definite matrix. Relate the factorization

$$
\begin{equation*}
T \widetilde{U}=\widetilde{L} \tag{1}
\end{equation*}
$$

to the standard $L D L^{*}$ factorization of $T$ to prove that (1) always exists and it is unique. Here $\widetilde{U}$ is a unit (i.e., with 1's on the main diagonal) upper triangular matrix, and $\widetilde{L}$ is a lower triangular matrix.
(b) Let $\langle\cdot, \cdot\rangle$ be an arbitrary inner product in the vector space $\Pi_{n}$ (of all polynomials whose degree does not exceed $n$ ). Let $T$ be a positive definite moment matrix, i.e., $T=$ $\left[\left\langle x^{i}, x^{j}\right\rangle\right]_{i, j=0}^{n}$. Let

$$
\begin{equation*}
u_{k}(x)=u_{0, k}+u_{1, k} x+u_{2, k} x^{2}+\ldots+u_{k-1, k} x^{k-1}+x^{k} . \tag{2}
\end{equation*}
$$

be the $k$-th orthogonal polynomial with respect to $\langle\cdot, \cdot\rangle$. Prove that the $k$-th column of the matrix $\widetilde{U}$ of (a) contains the coefficients of $u_{k}(x)$ as in

$$
\widetilde{U}=\left[\begin{array}{ccccccc}
1 & u_{0,1} & u_{0,2} & u_{0,3} & \cdots & \cdots & u_{0, n} \\
0 & 1 & u_{1,2} & u_{1,3} & \cdots & \cdots & u_{1, n} \\
0 & 0 & 1 & u_{2,3} & \cdots & \cdots & u_{2, n} \\
\vdots & & 0 & 1 & \cdots & \cdots & u_{3, n} \\
\vdots & & & \ddots & \ddots & & \vdots \\
\vdots & & & & \ddots & 1 & u_{n-1, n} \\
0 & & & \cdots & \cdots & 0 & 1
\end{array}\right]
$$

(c) Assuming now that the moment matrix $T$ has Toeplitz structure derive the so-called Levinson algorithm, that is, an algorithm to compute the columns of $\widetilde{U}$ based on the formula (deduce it) that relates the $k$-th column $u_{k}$ of $U$ to its "predecessor" $u_{k-1}$ ( $k=2,3, \ldots, n$ ).

Hint: Use the fact (no need to prove it) that Toeplitz moment matrices $T$ have the following property: if

$$
T\left[\begin{array}{c}
x_{1} \\
x_{2} \\
x_{3} \\
\vdots \\
x_{n-2} \\
x_{n-1} \\
x_{n}
\end{array}\right]=\left[\begin{array}{c}
y_{1} \\
y_{2} \\
y_{3} \\
\vdots \\
y_{n-2} \\
y_{n-1} \\
y_{n}
\end{array}\right]
$$

then

$$
T\left[\begin{array}{c}
x_{n}^{*} \\
x_{n-1}^{*} \\
x_{n-2}^{*} \\
\vdots \\
x_{3}^{*} \\
x_{2}^{*} \\
x_{1}^{*}
\end{array}\right]=\left[\begin{array}{c}
y_{n}^{*} \\
y_{n-1}^{*} \\
y_{n-2}^{*} \\
\vdots \\
y_{3}^{*} \\
y_{2}^{*} \\
y_{1}^{*}
\end{array}\right]
$$

(d) Prove that the algorithm of (c) uses $O\left(n^{2}\right)$ arithmetic operations.

Q4) Solve (a), (b), (c)
(a) Use the fact that each norm $\|\cdot\|$ on $\mathbb{C}^{n}$ is uniformly continuous (no need to prove the latter fact, just formulate it as a specific inequality) to prove the following theorem.
All norms on $\mathbb{C}^{n}$ are equivalent in the following sense. For each pair of norms $p_{1}(x)$ and $p_{2}(x)$ there are positive constants $m$ and $M$ satisfying

$$
m p_{2}(x) \leq p_{1}(x) \leq M p_{2}(x)
$$

for all $x$.
(b) Prove that if $F$ is an $n \times n$ matrix with $\|F\|<1$, then $(I+F)^{-1}$ exists and satisfies

$$
\left\|(I+F)^{-1}\right\| \leq \frac{1}{1-\|F\|}
$$

(c) Let $A$ be a nonsingular $n \times n$ matrix, $B=A(I+F),\|F\|<1$, and $x$ and $\Delta x$ be defined by

$$
A x=b, \quad B(x+\Delta x)=b .
$$

Use (b) to prove that

$$
\frac{\|\Delta x\|}{\|x\|} \leq \frac{\|F\|}{1-\|F\|}
$$

as well as

$$
\frac{\|\Delta x\|}{\|x\|} \leq \frac{\operatorname{cond}(A)}{1-\operatorname{cond}(A) \frac{\|B-A\|}{\|A\|}} \cdot \frac{\|B-A\|}{\|A\|}
$$

if

$$
\operatorname{cond}(A) \frac{\|B-A\|}{\|A\|}<1 .
$$

