INSTRUCTIONS: Solve three out of four questions. You do not have to prove results which you rely upon, just state them clearly.

Good luck!

- Q1) Solve 4 problems out of (a), (b), (c), (d). (e)
 - (a) Consider the following classical interpolation problem.

Given n + 1 support points

(x_i, f_i) i = 0,...,n;
(x_i ≠ x_j for i ≠ j).

Find a polynomial P(x) whose degree does not exceed n such that

P(x_i) = f_i, i = 0,...n.

Define the Vandermonde matrix, and then reformulate the above interpolation problem as a matrix problem of solving a linear system of equations with the Vandermonde coefficient matrix.

Use the condition

$$x_i \neq x_j$$
 for $i \neq j$,

to prove that the Vandermonde matrix is nonsingular.

Use the latter fact to prove that the classical interpolation problem stated above has a unique solution.

(b) Let $P_{i_0i_1...i_k}(x)$ be the (unique) polynomial that interpolates at points

$$(x_{i_m}, f_{i_m}) \qquad m = 0, \dots k.$$

Prove the Neville formula

$$P_{0,1,2,\dots,k}(x) = \frac{(x-x_j)P_{0,1,\dots,j-1,j+1,\dots,k}(x) - (x-x_i)P_{0,1,\dots,i-1,i+1,\dots,k}(x)}{x_i - x_j}$$

(c) Prove that there exists a unique coefficient $f_{i_0...i_k}$ such that

$$P_{i_0\dots i_k}(x) = P_{i_0\dots i_{k-1}} + f_{i_0\dots i_k}(x - x_{i_0})(x - x_{i_1})\cdots(x - x_{i_{k-1}}).$$

(d) Prove the recursion:

$$f_{i_0\dots i_k} = \frac{f_{i_1\dots i_k} - f_{i_0\dots i_{k-1}}}{x_{i_k} - x_{i_0}}$$

(e) Prove the following theorem (error in polynomial interpolation).

If the function f has an (n + 1)st derivative, then for every argument \bar{x} there exist a number ξ (in the smallest interval containing $x_{i_0}, x_{i_1}, \ldots, x_{i_n}, \bar{x}$), satisfying

$$f(\bar{x}) - P_{i_0, i_1, \dots, i_n}(x) = \frac{w(\bar{x})f^{(n+1)}(\xi)}{(n+1)!},$$

where

$$w(x) = (x - x_{i_0})(x - x_{i_1})\dots(x - x_{i_n})$$

- **Q2)** Solve (a), (b), (c)
 - (a) Use the fact that each norm || · || on Cⁿ is uniformly continuous (no need to prove the latter fact, just formulate it as a specific inequality) to prove the following theorem. All norms on Cⁿ are equivalent in the following sense. For each pair of norms p₁(x) and p₂(x) there are positive constants m and M satisfying

$$mp_2(x) \le p_1(x) \le Mp_2(x)$$

for all x.

(b) Prove that if F is an $n \times n$ matrix with ||F|| < 1, then $(I+F)^{-1}$ exists and satisfies

$$||(I+F)^{-1}|| \le \frac{1}{1-||F||}.$$

(c) Let A be a nonsingular $n \times n$ matrix, B = A(I + F), ||F|| < 1, and x and Δx be defined by

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$$Ax = b,$$
 $B(x + \Delta x) = b$

Use (b) to prove that

$$\frac{\|\Delta x\|}{\|x\|} \le \frac{\|F\|}{1 - \|F\|}$$

as well as

$$\frac{\|\Delta x\|}{\|x\|} \leq \frac{\operatorname{cond}(A)}{1 - \operatorname{cond}(A) \frac{\|B - A\|}{\|A\|}} \cdot \frac{\|B - A\|}{\|A\|}$$

if

$$cond(A) \frac{\|B - A\|}{\|A\|} < 1.$$

- **Q3)** Answer 3 out of 4 questions (a), (b), (c), (d).
 - (a) Define a Hankel matrix. Let H be an $n \times n$ positive definite Hankel matrix. Relate the factorization

$$H\widetilde{U} = \widetilde{L} \tag{1}$$

to the standard LDL^* factorization of H to prove that (1) always exists and it is unique. Here \widetilde{U} is a unit (i.e., with 1's on the main diagonal) upper triangular matrix, and \widetilde{L} is a lower triangular matrix. (b) Let $\langle \cdot, \cdot \rangle$ be an inner product in the vector space Π_n (of all polynomials whose degree does not exceed n). Let the above Hankel matrix H be a moment matrix, i.e., $H = [\langle x^i, x^j \rangle]_{i,j=0}^n$. Let

$$u_k(x) = u_{0,k} + u_{1,k}x + u_{2,k}x^2 + \dots + u_{k-1,k}x^{k-1} + x^k.$$
(2)

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be the k-th orthogonal polynomial with respect to $\langle \cdot, \cdot \rangle$. Prove that the k-th column of the matrix \widetilde{U} of (a) contains the coefficients of $u_k(x)$ as in

	1	$u_{0,1}$	$u_{0,2}$	$u_{0,3}$	• • •	• • •	$u_{0,n}$	
	0	1	$u_{1,2}$	$u_{1,3}$	• • •	• • •	$u_{1,n}$	
-	0	0	1	$u_{2,3}$	•••	•••	$u_{2,n}$	
$\widetilde{U} =$:		0	1			$u_{3,n}$.
	•			·	·		÷	
-	:				۰.	1	$u_{n-1,n}$	
	0				•••	0	1	

- (c) Derive a algorithm to compute the columns of \widetilde{U} based on the formula (deduce it) that relates the k-th column u_k of U to its two "predecessors" u_{k-2}, u_{k-1} (k = 3, ..., n).
- (d) Prove that the algorithm of (c) uses $O(n^2)$ arithmetic operations.

Γ.

Q4) Answer 4 out of 5 questions (a), (b), (c), (d), (e).

Derive a fast $O(n \log n)$ FFT-based algorithm for the polynomial multiplication problem, that is, given coefficients of two polynomials a(x), b(x), compute the coefficients of their product c(x) = a(x)b(x).

- (a) Prove that the above polynomial multiplication problem is equivalent to the problem of multiplying a lower triangular Toeplitz matrix by a vector.
- (b) Show how to "embed" a Toeplitz matrix into a circulant matrix, and justify the fact that the problem of (a) (that is, of multiplying a lower triangular Toeplitz matrix by a vector) can be solved via multiplying a circulant matrix by a vector.
- (c) Prove that any circulant matrix C admits a factorization

$$C = FDF^*$$

where F is the DFT matrix and D is a diagonal matrix.

- (d) Deduce the formula for the diagonal entries of D.
- (e) Describe "in words" how the results of (a), (b), (c), and (d) allow us to compute the coefficients of c(x) = a(x)b(x) in $O(n \log n)$ arithmetic operations.